

I. Short answer 1-12. Do not give proofs (3pts each).

1. Compute the following:

$$28 \operatorname{div} 3 = \underline{9} \quad 28 \operatorname{mod} 3 = \underline{1}$$

$$-17 \operatorname{div} 6 = \underline{-3} \quad -17 \operatorname{mod} 6 = \underline{1}$$

Division Algorithm

$$28 = 9 \cdot 3 + 1$$

$$-17 = -3 \cdot 6 + 1$$

$$(-17 = -2 \cdot 6 - 5, \text{ but } -5 < 0)$$

2. What condition on positive integers a and b must hold for them to be *relatively prime*?

$$\underline{\operatorname{gcd}(a, b) = 1}$$

3. Give an example of a function $f: \mathbb{R} \rightarrow \mathbb{R}$ that is one-to-one but not onto.

$$f(x) = e^x$$

4. Define $f: \mathbb{R} \rightarrow \mathbb{R}$ by $f(x) = x^3$, and $g: \mathbb{R} \rightarrow \mathbb{R}^+$ by $g(x) = 2^x$. Evaluate the following:

$$(g \circ f)(-1) = \underline{\frac{1}{2}}$$

$$(g \circ f)^{-1}(16) = \underline{\sqrt[3]{4}}$$

$$(g \circ f)(-1) = g(f(-1)) = g((-1)^3)$$

$$= g(-1) = 2^{-1} = \frac{1}{2}$$

$$(g \circ f)^{-1}(16) = (f^{-1} \circ g^{-1})(16) = f^{-1}(g^{-1}(16)) = f^{-1}(\log_2 16)$$

$$= f^{-1}(4) = \sqrt[3]{4}$$

5. The first and second terms of an arithmetic progression are $a_1 = 5.5$ and $a_2 = 8$. What is a_4 ?

$$a_4 = \underline{13}$$

$$a_1 = 5.5 \quad a_2 = 8 \quad a_3 = 10.5 \quad a_4 = 13$$

6. Assume for two positive integers a and b that $a \cdot b = 2^4 \cdot 3^3 \cdot 5 \cdot 7^2$ and $\operatorname{gcd}(a, b) = 2^2 \cdot 3 \cdot 7$. What is $\operatorname{lcm}(a, b)$?

$$\underline{2^2 \cdot 3^2 \cdot 5 \cdot 7}$$

7. Compute the following:

$$\operatorname{gcd}(0, 4) = \underline{4} \quad \operatorname{gcd}(-12, -50) = \underline{2}$$

$$12 = 2^2 \cdot 3 \quad 50 = 2 \cdot 5^2$$

8. Order the following expressions in terms of **increasing** growth rate: n^2 , $n!$, $n \log_2 n$, 2^n , $\log_2 n$

$\log_2 n$ $n \log_2 n$ n^2 2^n $n!$

9. In the best big-oh notation, how many multiplications are needed for Horner's method to evaluate a polynomial of degree n at a constant c ?

$O(n)$

Horner's method

procedure *Horner*(c, a_0, a_1, \dots, a_n : real numbers)

$y := a_n$

for $i := 1$ **to** n

$y := y * c + a_{n-i}$

10. Name an algorithm we studied for which the best-case time complexity is different than the worst-case time complexity (in terms of the representative operation we discussed).

linear search
(also binary search)

11. Circle the prime numbers and only the prime numbers in this list:

73, 92, 93, 109, 127, 133 2,3,5,7 ~~73~~ 2|92 3|93
2,3,5,7 ~~109~~ 2,3,5,7,11 ~~127~~ 7|133

12. A 2 cent and a 5 cent stamp are available (without limit) to make postage. What is the smallest amount of postage N in cents such that every $N, N+1, N+2, N+3, \dots$ can be composed of 2 cent and 5 cent stamps?

$N =$ 4 1 no 4 yes 8
2 yes 5 yes \vdots
3 no 6 = 4+2
7 = 5+2

II. Computation Problems 13-16 (10pts each). For full credit, show work to clearly justifying your answer.

13. Use the Euclidean Algorithm to express $\gcd(222, 180)$ as $r \cdot 222 + s \cdot 180$ for some $r, s \in \mathbb{Z}$.

$$\begin{aligned} 222 &= 1 \cdot 180 + 42 \\ 180 &= 4 \cdot 42 + 12 \\ 42 &= 3 \cdot 12 + \boxed{6} \\ 12 &= 2 \cdot 6 + 0 \end{aligned}$$

$$\begin{aligned} \gcd(222, 180) &= 6 \\ 6 &= 42 - 3 \cdot 12 \\ 6 &= 42 - 3(180 - 4 \cdot 42) \\ 6 &= 13 \cdot 42 - 3 \cdot 180 \\ 6 &= 13(222 - 1 \cdot 180) - 3 \cdot 180 \\ 6 &= \underbrace{13 \cdot 222} - \underbrace{16 \cdot 180} \\ \boxed{r=13} \quad \boxed{s=-16} \end{aligned}$$

14. Compute the sum $\sum_{i=0}^{10} (2 \cdot 3^{i+2} - 4 \cdot i)$ (an unsimplified answer is acceptable but must not contain ellipses (\dots) or i).

$$\begin{aligned} & \sum_{i=0}^{10} (2 \cdot 3^2 \cdot 3^i - 4 \cdot i) \\ &= 12 \sum_{i=0}^{10} 3^i - 4 \sum_{i=0}^{10} i \\ &= \boxed{12 \frac{3^{11}-1}{3-1} - 4 \frac{(10)(11)}{2}} \end{aligned}$$

15. Compute $(60032 \cdot 24005 + 90511 \cdot 3030) \pmod 3$. Shortcuts are recommended, but must be clear from your work.

$$\begin{aligned} &= (60032 \pmod 3)(24005 \pmod 3) + (90511 \pmod 3)(3030 \pmod 3) \\ & \pmod 3 \\ &= 2 \cdot 2 + 1 \cdot 0 \pmod 3 \\ &= 4 \pmod 3 \\ &= \boxed{1} \end{aligned}$$

16. Let $f(x) = 5x^2 + 10 \log_2 x$. Compute witnesses C and k that show $f(x)$ is $O(x^2)$. (Hint: $x > k \rightarrow |f(x)| \leq C|x^2|$.)

Pick
 $k=1$.

$$5|x^2| \leq 5|x^2| \text{ for } x > 1.$$

$$10|\log_2 x| \leq 10|x^2| \text{ for } x > 1.$$

$$\text{with } C=15, \quad x > 1 \rightarrow$$

$$\begin{aligned} 15x^2 + 10 \log_2 x &\leq 5|x^2| + 10|\log_2 x| \\ &\leq 5|x^2| + 10|x^2| \\ &= 15|x^2| = C|x^2|. \end{aligned}$$

$$\boxed{k=1, C=15}$$

III. Proofs 17-18 (12pts each). Partial credit for good proof structure.

17. Prove the following. Let m, n be positive integers greater than 1, and let a, b be integers. If $n|m$ and $a \equiv b \pmod{m}$, then $a \equiv b \pmod{n}$.

(Direct Proof)

Assume $n|m$ and $a \equiv b \pmod{m}$.

Since $n|m$, $n \cdot k = m$ for some $k \in \mathbb{Z}$.

Since $a \equiv b \pmod{m}$, $m|a-b$ by definition of congruence mod m .

Then $m = d = a - b$ for some $d \in \mathbb{Z}$.
substituting for m ,

$$n \cdot k \cdot d = a - b.$$

$k \cdot d$ is an integer since $k, d \in \mathbb{Z}$.

By definition of congruence mod m ,

$$a \equiv b \pmod{n}. \quad \square$$

18. Prove by a careful induction argument that every positive integer can be written in the form $2^r \cdot m$, where r is an integer and m is an odd integer.

For all $k \in \mathbb{Z}^+$, define $P(k)$ to be the statement

$$k = 2^r \cdot m \quad \text{for some } r \in \mathbb{Z} \text{ and odd } m \in \mathbb{Z}.$$

Proof by strong induction

Base case $k=1$ $1 = 2^0 \cdot 1$, and $0 \in \mathbb{Z}$, $1 \in \mathbb{Z}$ is odd.

Inductive step Let $k \in \mathbb{Z}^+$ and assume $P(1), \dots, P(k)$ all are true.

Consider $k+1$.

case 1 $k+1$ is odd.

Then $k+1 = 2^0(k+1)$, and $0 \in \mathbb{Z}$, $k+1 \in \mathbb{Z}$ is odd.

case 2 $k+1$ is even.

By definition $k+1 = 2h$ for $h \in \mathbb{Z}$, and

$$k+1 > 0 \rightarrow h > 0, \text{ and } 1 \leq h \leq k.$$

Since $P(h)$ is true,

$$h = 2^r \cdot m, \quad \text{where } r \in \mathbb{Z} \text{ and } m \in \mathbb{Z} \text{ is odd.}$$

Then $k+1 = 2h = 2^{r+1} \cdot m$, where $r+1 \in \mathbb{Z}$ and $m \in \mathbb{Z}$ is odd.

Therefore $P(k+1)$ is true.

By strong induction, $\forall n \in \mathbb{Z}^+$ $P(n)$ is true. \square